

Efficiency and Risk Trade-offs in Dynamic Oligopoly Markets

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- [1] Q. Huang, M. Roozbehani, M. Dahleh, "Efficiency and Risk Impacts of Market Architecture in Power Systems under Real-Time Pricing", To Appear in IEEE CDC, 2012
- [2] ____, "Efficiency and Risk Trade-offs in Dynamic Oligopoly Markets", work in progress

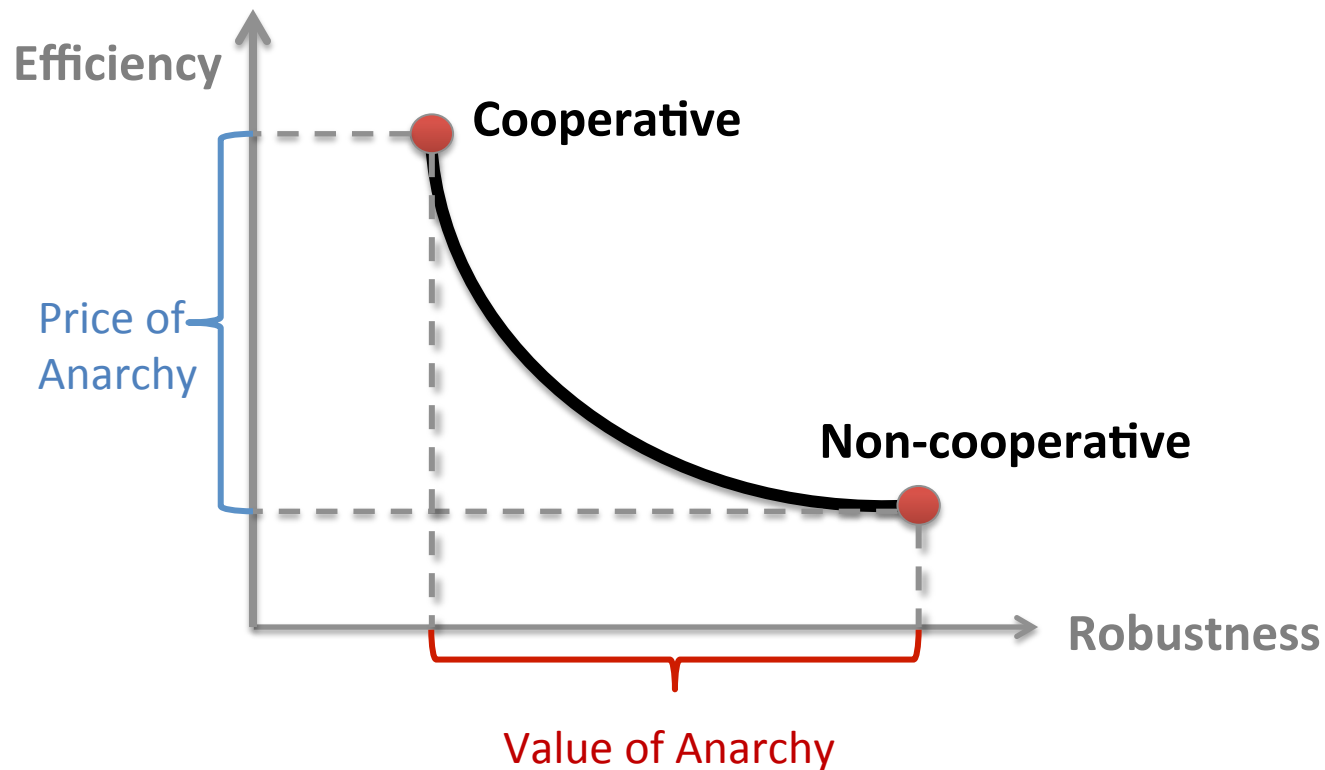
Exogenous shocks VS endogenous risk
Tail risk, aggregate demand / price spikes

- How does **endogenous risk** emerge in multi-agent dynamic system?
- What's the impact of local incentives on system-wide efficiency and risk?

market architectures → agent behaviors → aggregate outcome

Efficiency VS Risk

Tradeoff



cooperation Reduce externality, more efficient
More vulnerable to some exogenous shocks

system dynamics, uncertainties, frictions (deadline constraints)



- Previous works on **endogenous risk**
 - Heterogeneous beliefs [1]
 - Failure of the agents to rationalize feedback links [2]
- Our work
 - Agents are fully aware of pricing mechanism, have perfect information of system state, form rational expectations about other agents in the market.
 - Market architecture → Endogenous risk
 - Tradeoff
- Approach
 - Case study : Dynamic stochastic game, MDP
 - LTI reformulation and tradeoff analysis

[1] The leverage cycle (Geanakoplos 2009)

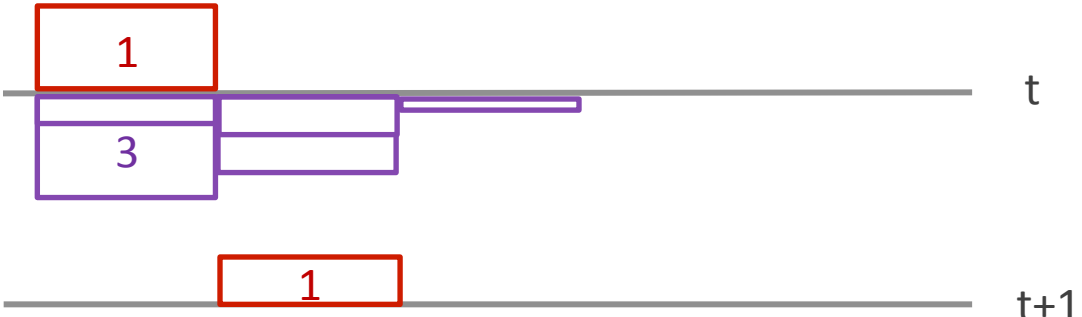
[2] "Endogenous Risk" (Danielsson and Shin 2003)

Setup

Agent arrival

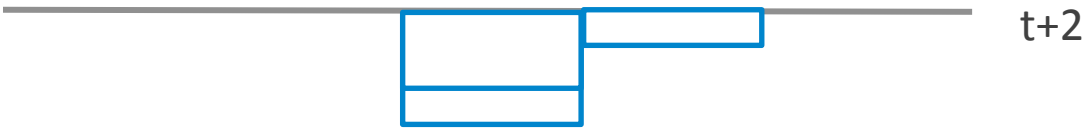
- **L** types (deadline constraint)
- Uncertainty
 - Bernoulli arrival
 - Workload distribution

$$d_l(t), l \in \{1, \dots, L\}$$



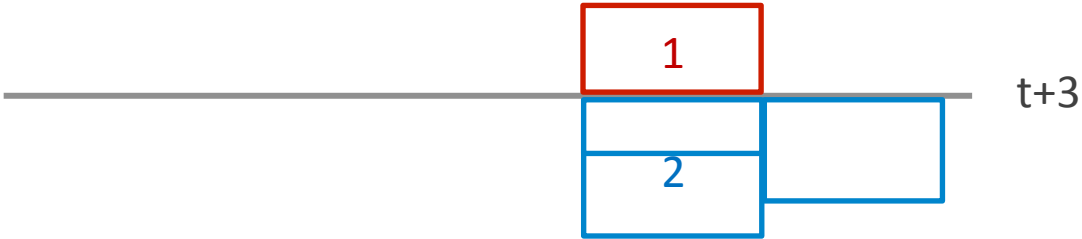
Decision

- State information
- Load scheduling



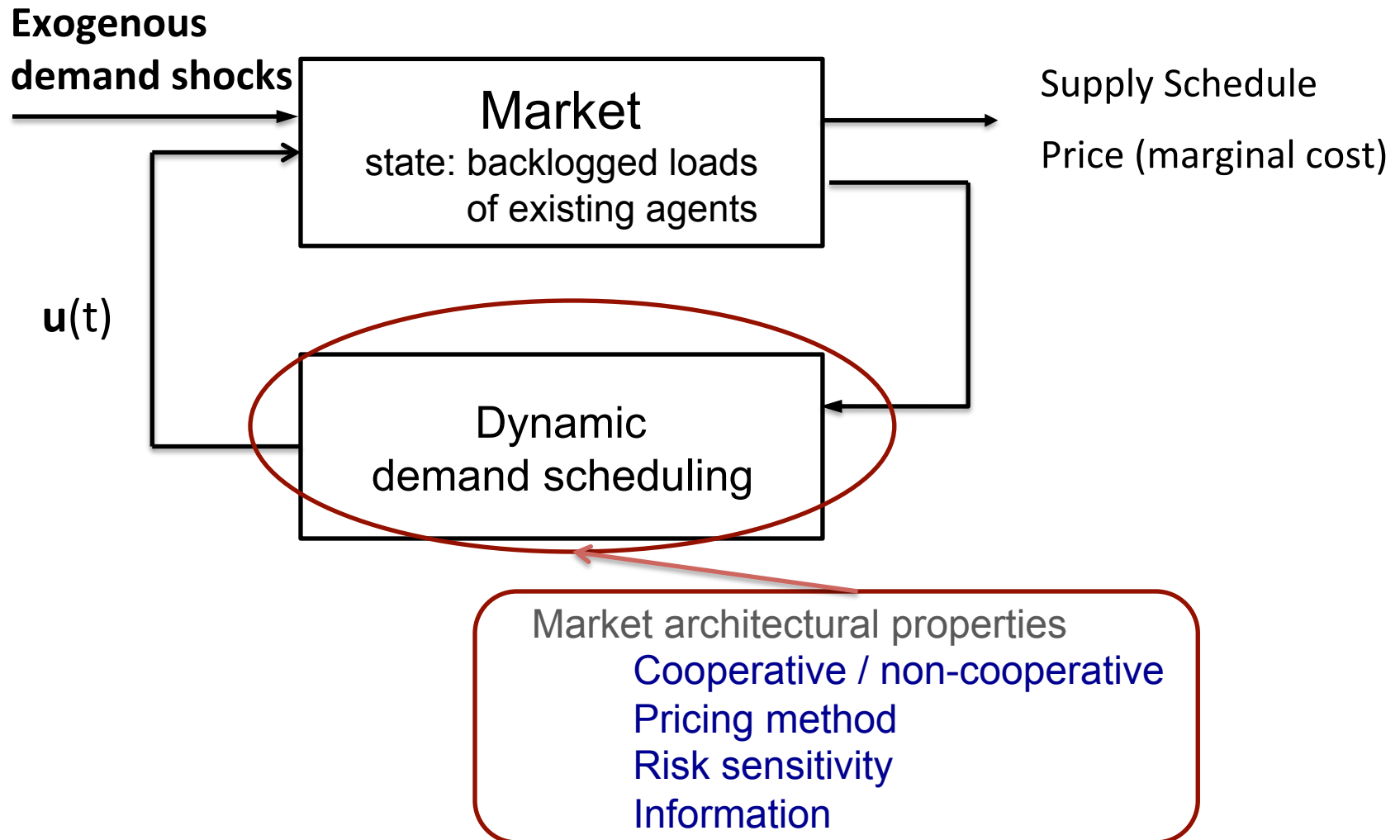
Cost $\mathbf{E}[\sum_{\tau=t}^{t+l} u_i(\tau)p(\tau)]$

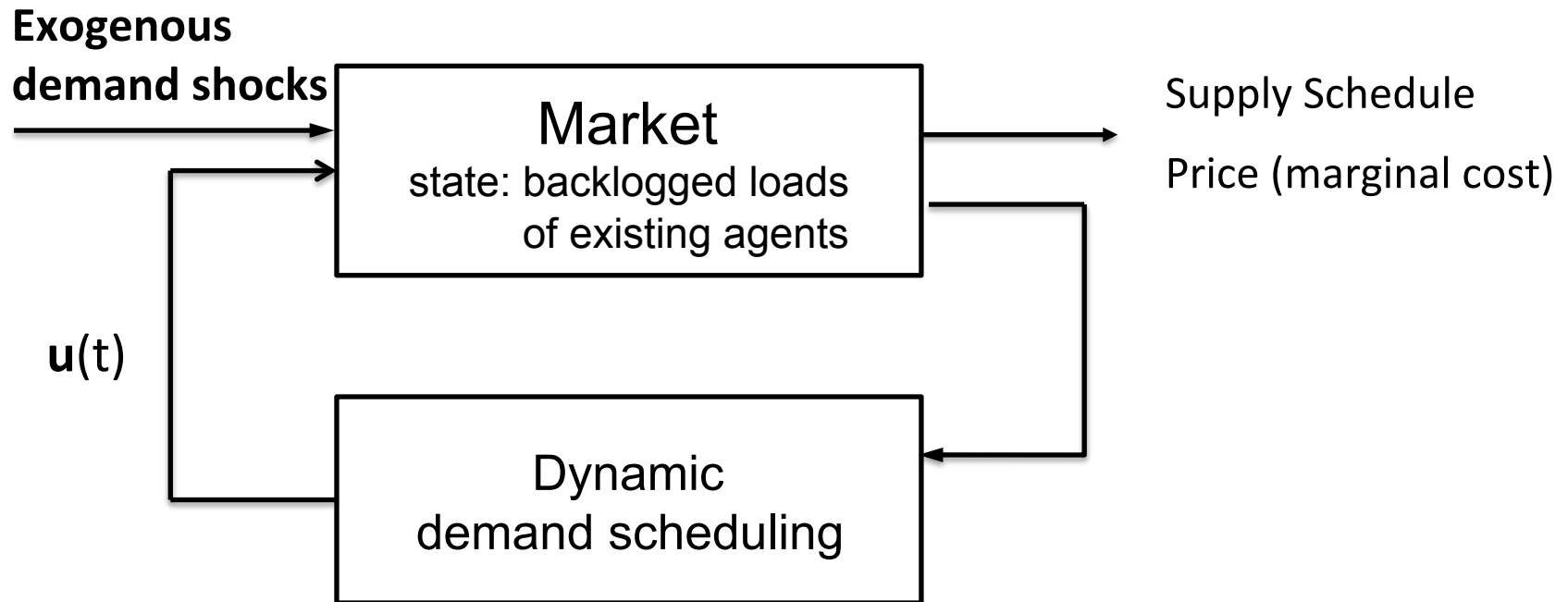
Pricing $p(t) = \sum_i u_i(t)$



System performance

- **Efficiency:** expected average cost (variance of aggregate demand process)
- **Risk:** aggregate demand spikes (tail probability of aggregate demand process)





Applications:

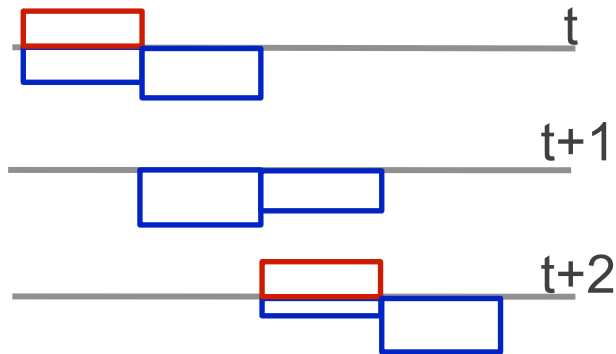
- consumer response to real-time pricing in power grids
- load scheduling in cloud computing market
- multi-portfolio execution problem
- consumption risk sharing

Case study L=2

Solutions



market architectures \rightarrow agent behaviors \Rightarrow aggregate demand process



- At most 1 decision maker at each period
 $\left(u(t), d_2(t) - u(t)\right)$
- Linear quadratic

$$u^s(x, d_2) = -a^s x + b^s d_2 + e^s$$

Non-cooperative

Dynamic stochastic game
Markov perfect equilibrium

Cooperative

Infinite horizon average cost MDP

$$a^{nc} < a^c \quad b^{nc} > b^c$$

In cooperative scheme, agents respond more aggressively to other agents' shock

Low efficiency
Low risk

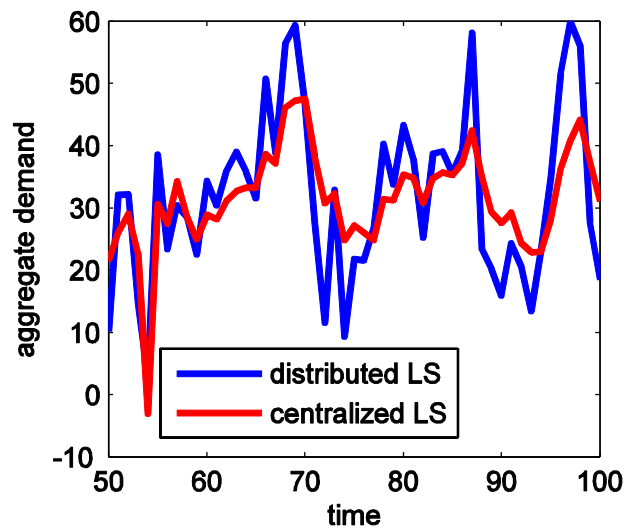
High efficiency
High risk

Case study L=2

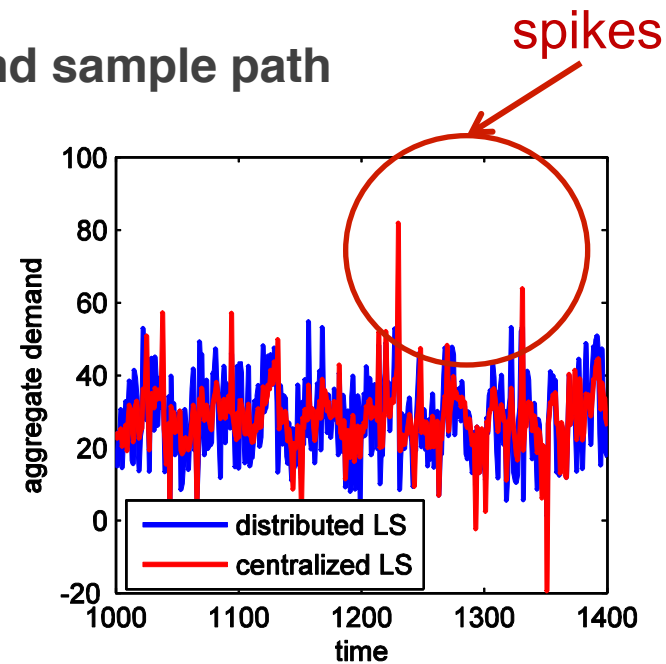
Results



Aggregate demand sample path



small time scale



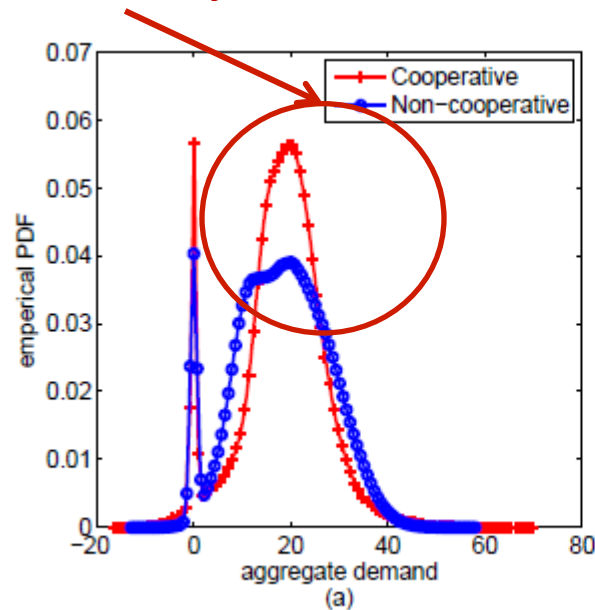
large time scale



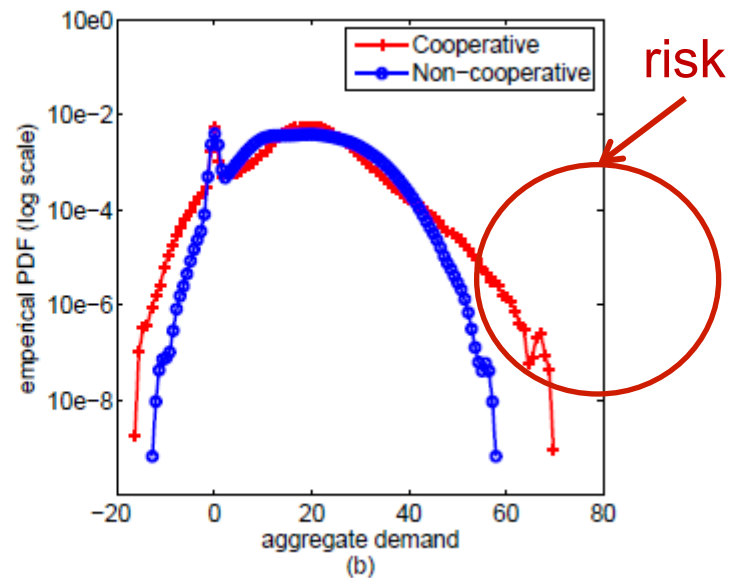
Cooperative
Non-cooperative

Aggregate demand distribution

efficiency



Linear scale



Log scale



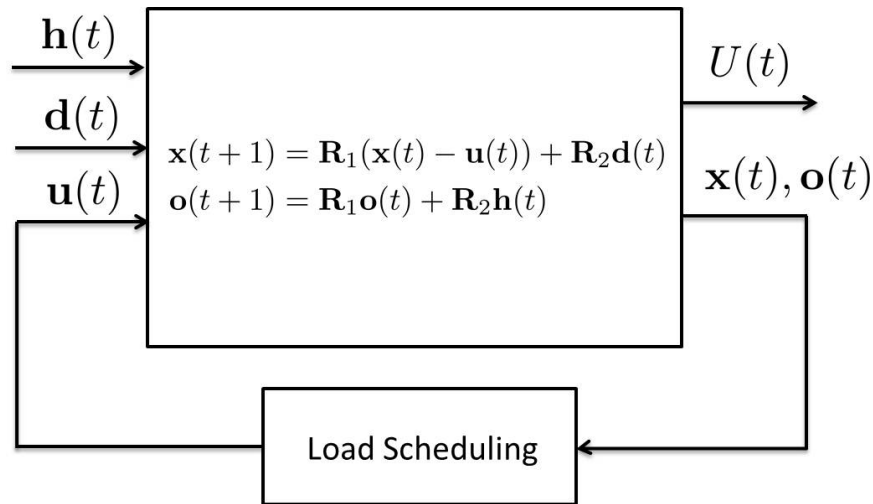
Cooperative
Non-cooperative

General L analysis

Modified system dynamics



- Jump linear system
- No closed form solution



(a)

Risk of spikes

High backlog of load

Absence of flexible load (Bernoulli)



Cooperative load scheduling

less externality

more rely on each other

more heavily use of backlog

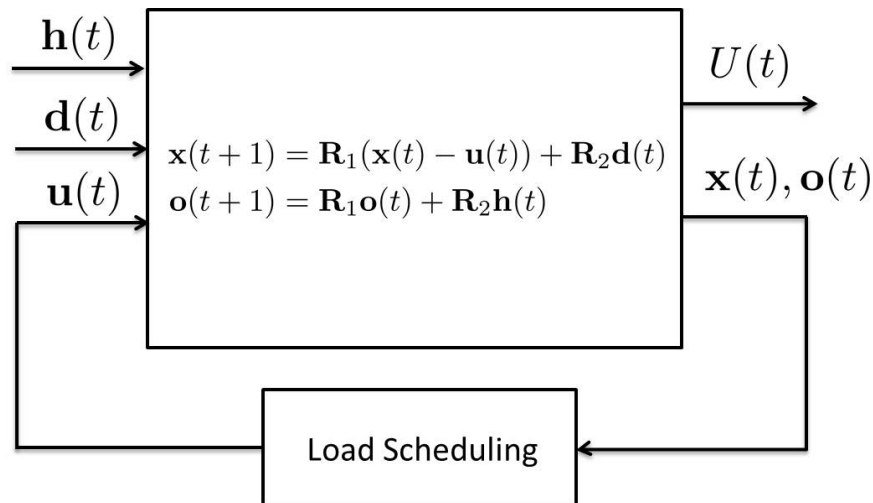
General L analysis

Modified system dynamics

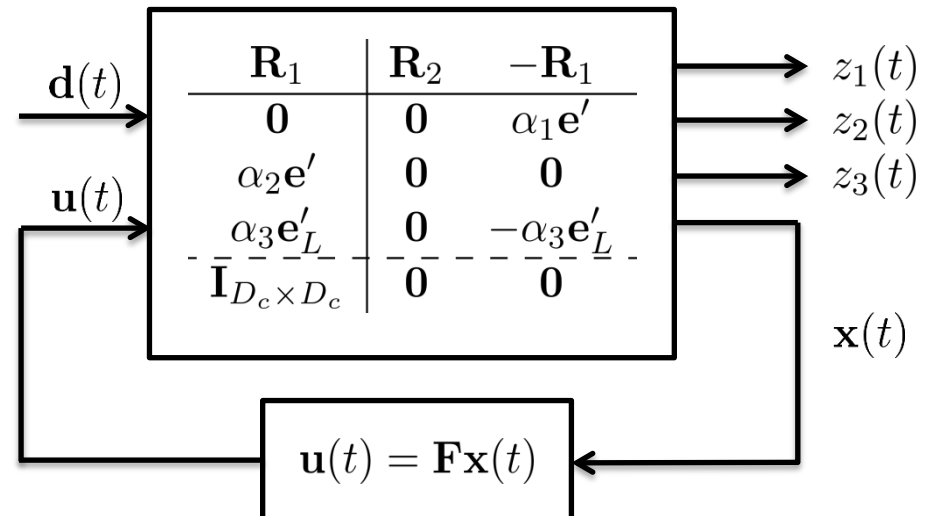


- Focus on linear dynamics
- Relax deadline constraints
- Substitute risk measure

- LTI system
- Multi-objective optimization
 - **Efficiency** aggregate demand 2nd moment
 - **Risk** aggregate backlog 2nd moment
 - **Aggregate unsatisfied load** 2nd moment



(a)



(b)

Efficiency frontier characterization (optimal **F** design with weighted outputs)

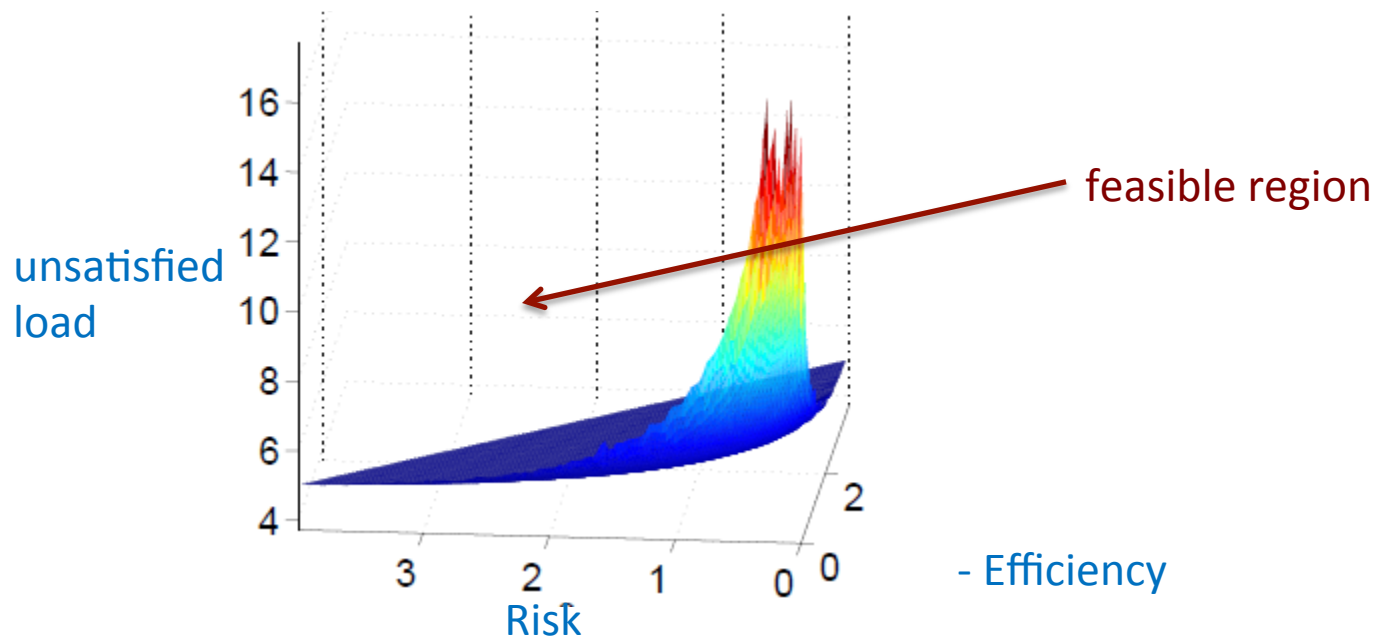
Three way tradeoff:

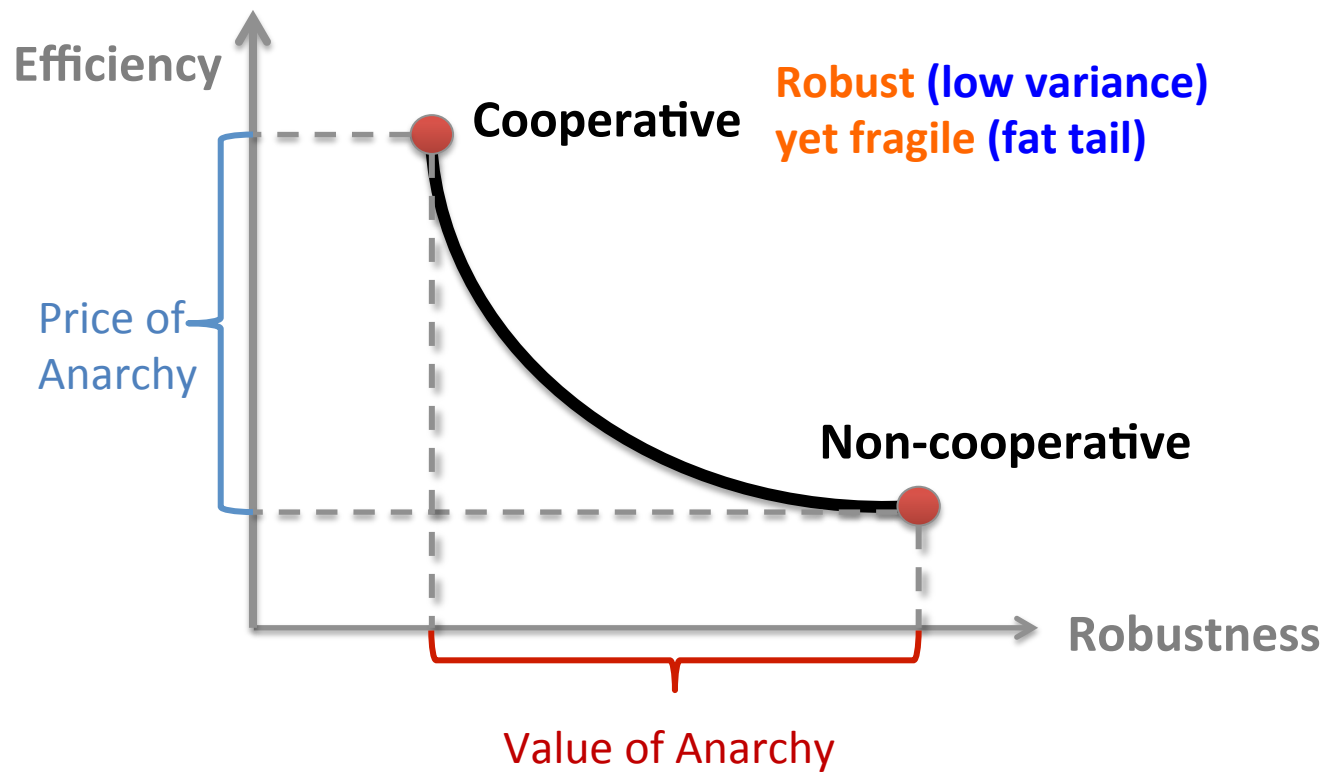
- - Efficiency
- Risk
- Unsatisfied load

aggregate demand 2nd moment

aggregate backlog 2nd moment

aggregate unsatisfied load 2nd moment





Efficiency

System
design

Dynamic pricing rule?

**Outcome in
multi-agent
systems**

Robustness

Other goals
(e.g. complexity...)



Thank you

System state: $s(t) = (x(t), d_2(t))$

- Aggregate unshiftable loads $x(t)$
- Consumer arrival with flexible load $d_2(t)$

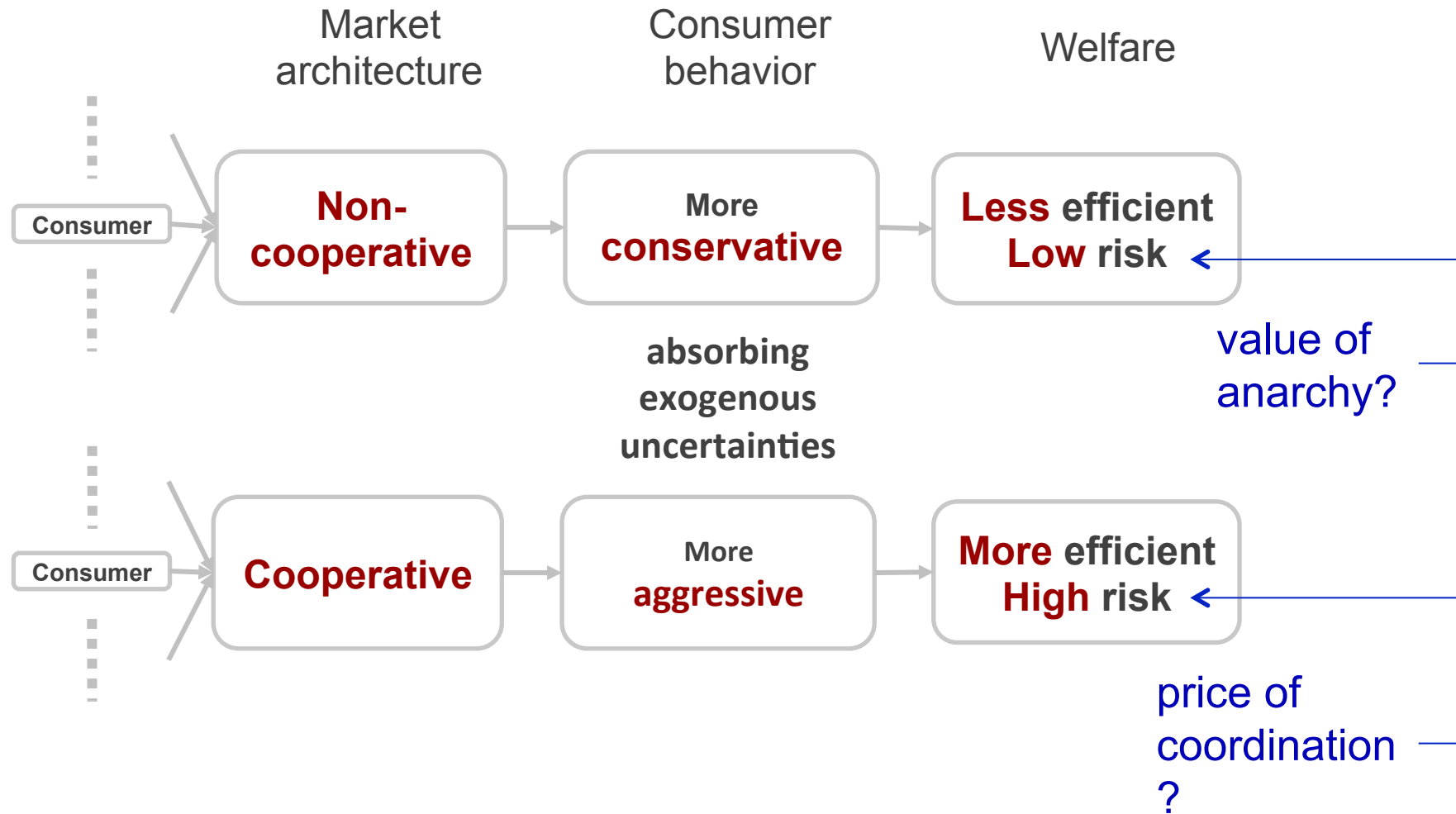
$$\underbrace{x(t)}_{\text{aggregate unshiftable}} = \underbrace{d_1(t)}_{\text{unshiftable arrival at current period}} + \underbrace{d_2(t-1) - u(t-1)}_{\text{leftover from last period's shiftable}}$$

Load scheduling decision:

- At most 1 decision maker at t : the new type 2 agent
- Split load into two periods $(t, t+1)$ based on state information

$$(u(t), d_2(t) - u(t))$$

Efficiency and Risk Implications of Architecture





market architectures \rightarrow agent behaviors \rightarrow aggregate demand process

Efficiency

Producer surplus + Consumer surplus **Variance**

$$W = \underbrace{\mathbf{E}[p(t)U(t) - \frac{1}{2}U(t)^2]}_{W_p} + \underbrace{\mathbf{E}[-p(t)U(t)]}_{W_c} = -\frac{1}{2}\mathbf{E}[U(t)^2]$$

Risk

Probability of aggregate unshiftable load peak **Tail probability**

$$\Pr(x(t) \geq M)$$

Case study L=2

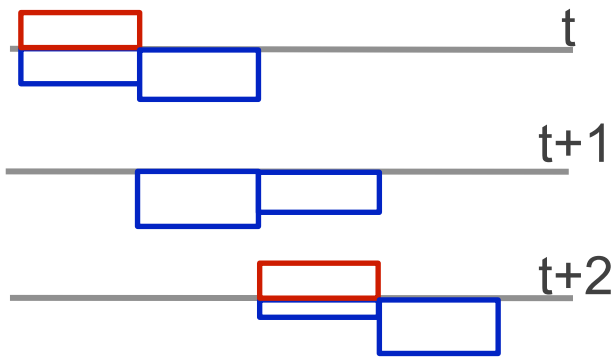
Solution: Non-Cooperative Case



- Full state (and everything else) information
- No coordination among strategic agents
- Focus on steady state in symmetric equilibria

Symmetric Markov Perfect equilibrium in dynamic stochastic game

$$u^s(x(t), d_2(t)) = \arg \min_u \{p(t)u + \mathbf{E}_t[p(t+1)(d_2(t) - u)]\}$$



$$p(t) = x(t) + u$$
$$p(t+1) = x(t+1) + u^s(x(t+1), d_2(t+1))$$

Overlapping type 2 consumers

Flavor of Stackelberg competition

Symmetric Markov Perfect equilibrium in dynamic stochastic game

$$u^s(x(t), d_2(t)) = \arg \min_u \{p(t)u + \mathbf{E}_t[p(t+1)(d_2(t) - u)]\}$$

Equilibrium strategy

There exists a unique symmetric MPE with linear equilibrium strategy:

$$u^s(x, d_2) = -a^s x + b^s d_2 + e^s$$

where a^s, b^s, e^s are constants determined by q_1, q_2, μ_1, μ_2 .

Case study L=2

Solution: Cooperative Case



Bellman equation for infinite horizon average cost MDP

$$\lambda^c + V^c(x) = (1 - q_2)(x^2 + \mathbf{E}[V^c(d_1)]) + q_2 \mathbf{E}[\min_u \{(x + u)^2 + V^c(d_2 - u + d_1)\}]$$

Optimal stationary policy

There exists an optimal stationary load scheduling policy:

$$u^c(x, d_2) = -a^c x + b^c d_2 + e^c$$

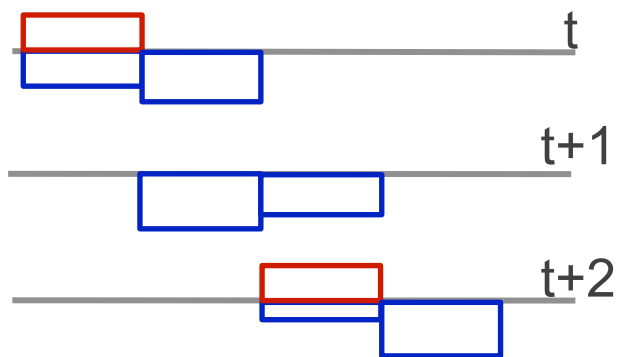
where a^c, b^c, e^c are constants determined by q_1, q_2, μ_1, μ_2 .

Case study L=2

Solutions



market architectures \rightarrow agent behaviors \rightarrow aggregate demand process



System performance

Efficiency: expected average cost (variance)

Risk: aggregate demand spikes (tail probability)

Non-cooperative

Cooperative

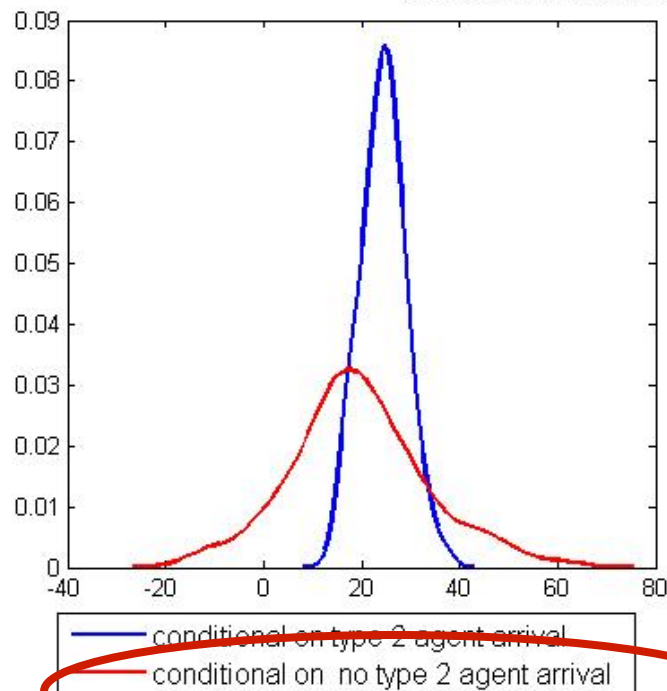
$$A_{nc} > a_c, b_{nc} < b_c$$

Low efficiency
Low risk

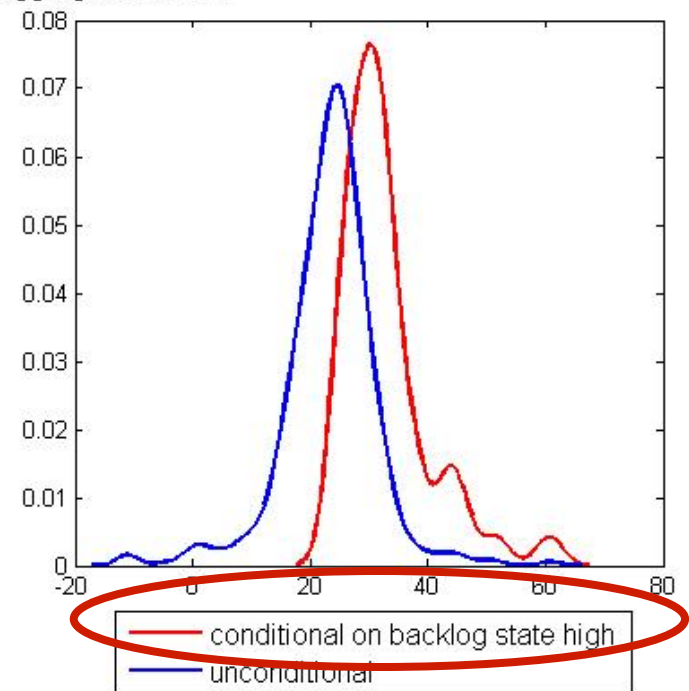
High efficiency
High risk

When do spikes occur?

conditional distribution of the aggregate demand



Absence of flexible load



High backlog state